

Quantum Information Measures in QFT

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Quantum Information

Measures in QFT

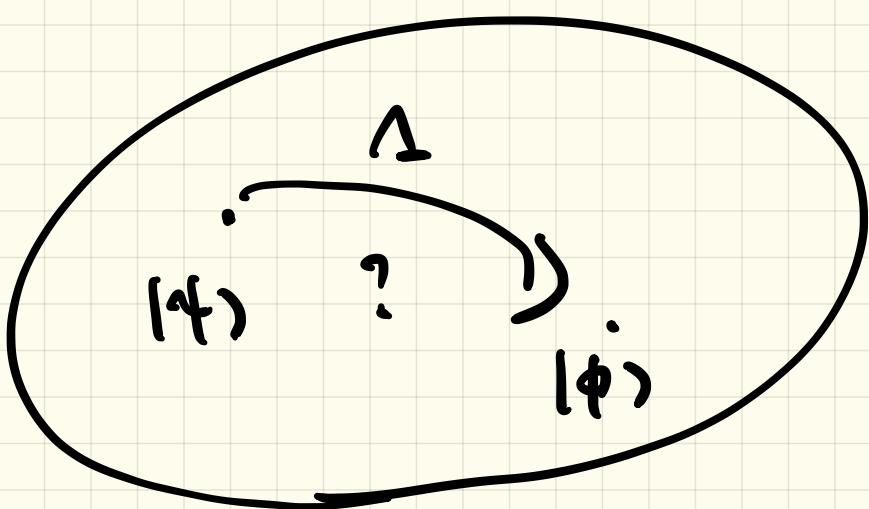
0. Introduction

III Quantum Information Theory
(QIT)

state conversion

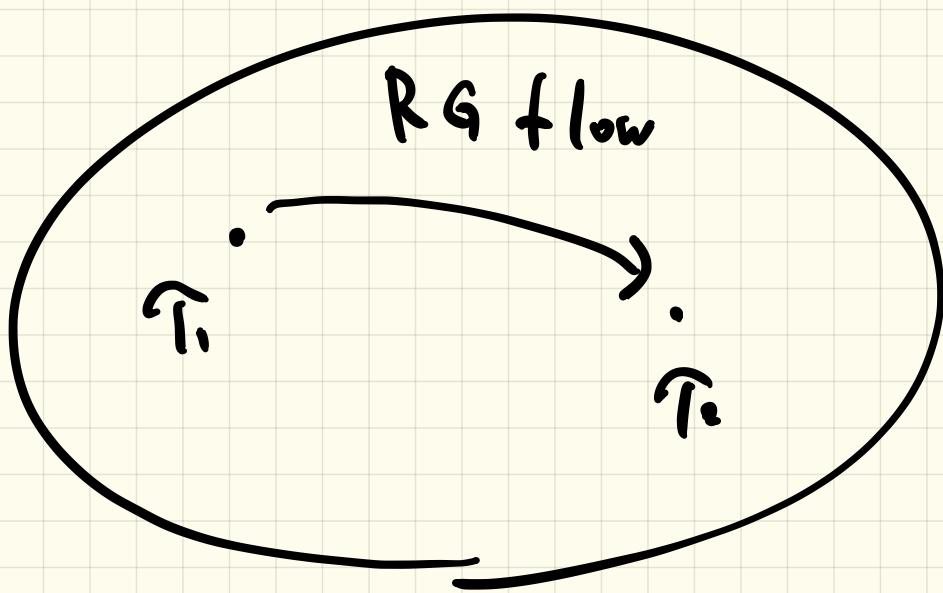
$$|\psi\rangle \xrightarrow{A} |\phi\rangle$$

- What operations Λ are allowed?
- When $|\psi\rangle$ is transformed to $|\phi\rangle$?
 \rightsquigarrow Ordering of quantum states



Space of quantum states

(6) QFT



Space of QFTs

Under RG flow

$$\widehat{T_1} \longrightarrow \widehat{T_2}$$

- When $\widehat{T_1}$ flows to $\widehat{T_2}$?

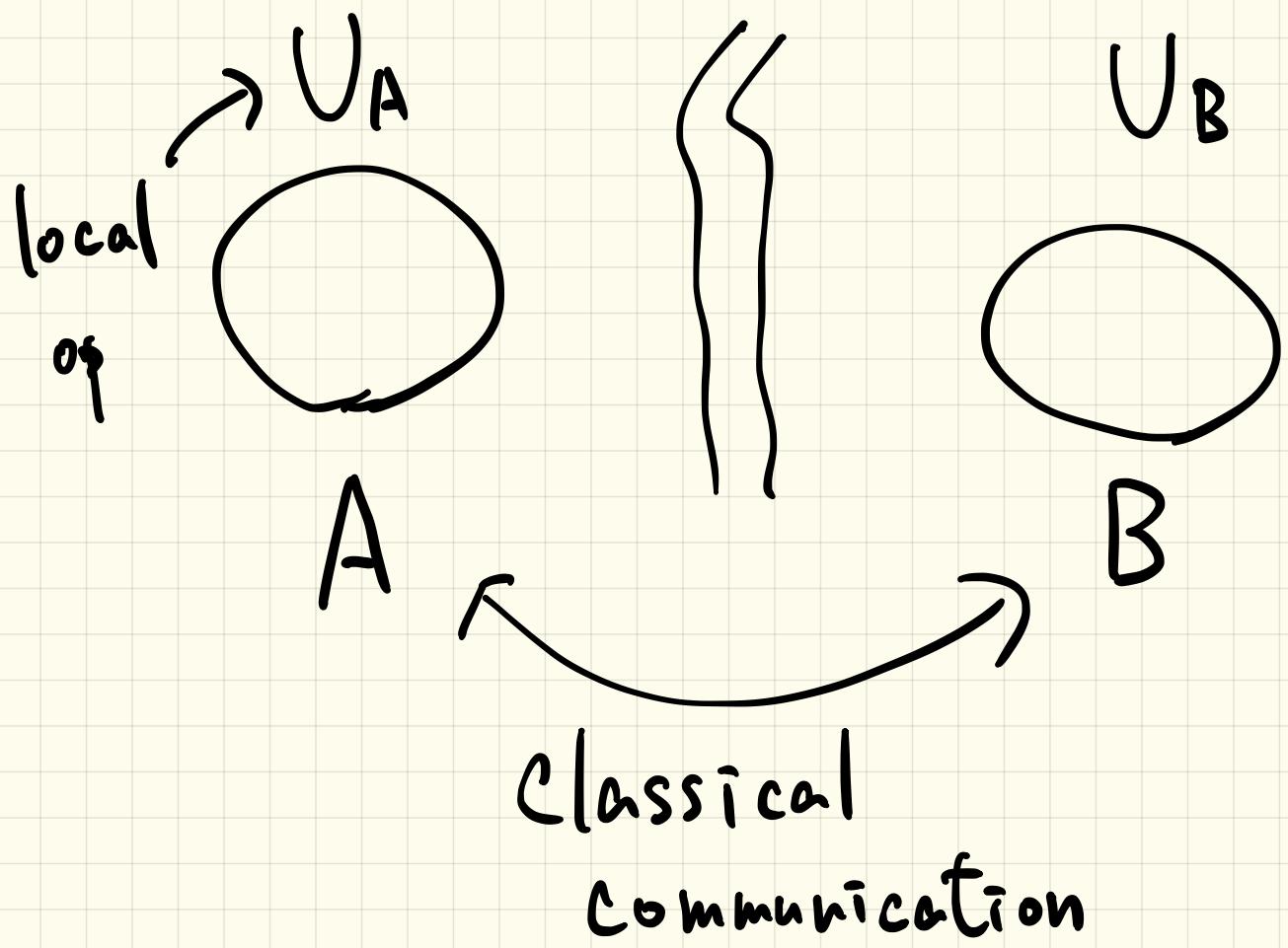
→ Ordering QFTs
(C-theorems)

1. QIT

|ψ⟩ : quantum states

Λ : LOCC

local op classical
communication



Classification of quantum states

- pure bipartite system

$$\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_B$$

$$|\psi\rangle \leftrightarrow P_\psi = |\psi\rangle\langle\psi|$$

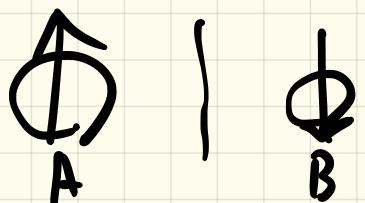
(mixed state)

$$\leftrightarrow \rho = \sum_i p_i |\psi_i\rangle\langle\psi_i|$$

- Separable state

$$|\psi\rangle = |\psi_A\rangle \otimes |\psi_B\rangle$$

\uparrow \uparrow
 \mathcal{H}_A \mathcal{H}_B



- Entangled state

= non-separable

Bell pair (EPR)

$$|\Phi^+\rangle = \frac{1}{\sqrt{2}}(|0_A\rangle|1_B\rangle + |1_A\rangle|0_B\rangle)$$
$$\neq |\psi_A\rangle \otimes |\psi_B\rangle$$

for any ψ_A, ψ_B

Schmidt decomposition

Canonical form

of a bipartite quantum state

In general

$$|\psi\rangle = \sum_{i,j} c_{ij} |i_A\rangle |j_B\rangle$$

\Downarrow

\mathcal{H}_A \mathcal{H}_B

CPX

By changing the orthonormal basis

$$|\tilde{i}_A\rangle \rightarrow V_{ik} |k_A\rangle \quad \left| \begin{array}{l} U, V \\ \text{unitaries} \end{array} \right.$$
$$|\tilde{j}_B\rangle \rightarrow V_{jl} |l_B\rangle$$

C_{ij} : $d_A \times d_B$ rectangular

$$C_{ij} = \left(\begin{array}{c} \xleftarrow{d_A} & \xrightarrow{d_B} \\ \text{CPX matrix} \end{array} \right)$$

Singular value decomposition

$$C_{ij} = (U \cdot \Lambda \cdot V)_{ij}$$

$$\Lambda = \begin{pmatrix} \sqrt{\lambda_1} & & 0 & \\ & \ddots & & 0 \\ 0 & & \sqrt{\lambda_{d_B}} & \\ & & & 0 \end{pmatrix}$$

$\xrightarrow{d_A}$ $\downarrow d_B$

$$(d_A \geq d_B)$$

- $\lambda_i \geq 0$
- $(\Psi | \Psi) = \left| \sum_{i=1}^{\min(d_A, d_B)} \sqrt{\lambda_i} |i_A\rangle |i_B\rangle \right|^2$
- $(\Psi | \Psi) = \left| \sum_{i=1}^{\min(d_A, d_B)} \lambda_i \right|^2 = 1$
- $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_{\min(d_A, d_B)} \geq 0$

$$\lambda^+ = (\lambda_1^+, \lambda_2^+, \dots, \lambda_{\min(d_A, d_B)}^+)$$

↑
Schmidt vector

- We want to introduce "order" between quantum states

Partial order

For $x = (x_1, x_2, \dots)$

$y = (y_1, y_2, \dots)$



$$\left. \begin{array}{l} \in \mathbb{R}^n \\ \sum_{i=1}^n x_i \\ = \sum_{i=1}^n y_i \end{array} \right\}$$

$$x \prec y \quad \text{iff} \quad \sum_{i=1}^k x_i^\downarrow \leq \sum_{i=1}^k y_i^\downarrow$$

x is majorized

by y

for any k

Example

$$\begin{aligned} \left(\frac{1}{n}, \frac{1}{n}, -\frac{1}{n} \right) &\prec \left(\frac{1}{n-1}, \frac{1}{n-1}, -\frac{1}{n-1}, 0 \right) \\ &\prec \left(\frac{1}{n-2}, \frac{1}{n-2}, -\frac{1}{n-2}, 0, 0 \right) \\ &\prec \dots \\ &\prec (1, 0, \dots, 0) \end{aligned}$$

There are vectors
which cannot be compared

$$x \neq y, y \neq x$$

e.g.)

$$\begin{aligned} x &= \left(\frac{1}{2}, \frac{1}{4}, \frac{1}{4} \right) \\ y &= \left(\frac{2}{5}, \frac{2}{5}, \frac{1}{5} \right) \end{aligned} \quad \left| \begin{array}{l} k=1 \quad x_1 = \frac{1}{2} \Rightarrow \frac{2}{5} < y_1 \\ k=2 \\ x_1 + x_2 = \frac{3}{4} \\ y_1 + y_2 = \frac{4}{5} \end{array} \right.$$

QIT problem

When can a state $|\psi\rangle$

be transformed

into a state $|\phi\rangle$

by LOCC?

$$|\psi\rangle \xrightarrow{?} |\phi\rangle$$

LOCC \subset SEP (Separable op)

$$\Lambda(\rho) = \sum_m (A_m \otimes B_m) \rho (A_m \otimes B_m)^+$$

$$\boxed{\sum_m A_m^+ A_m \otimes B_m^+ B_m = \mathbb{1}_A \otimes \mathbb{1}_B}$$

LOCC

A performs op. A_i

A sends the results i to B

Then B performs some op

$$B_j^{(i)} \quad | \quad A_i \otimes B_j^{(i)}$$
$$m = (i, j)$$

Theorem ^{Michael} \rightarrow (Nielsen's majorization)

For bipartite pure states $| \psi \rangle, |\phi \rangle$

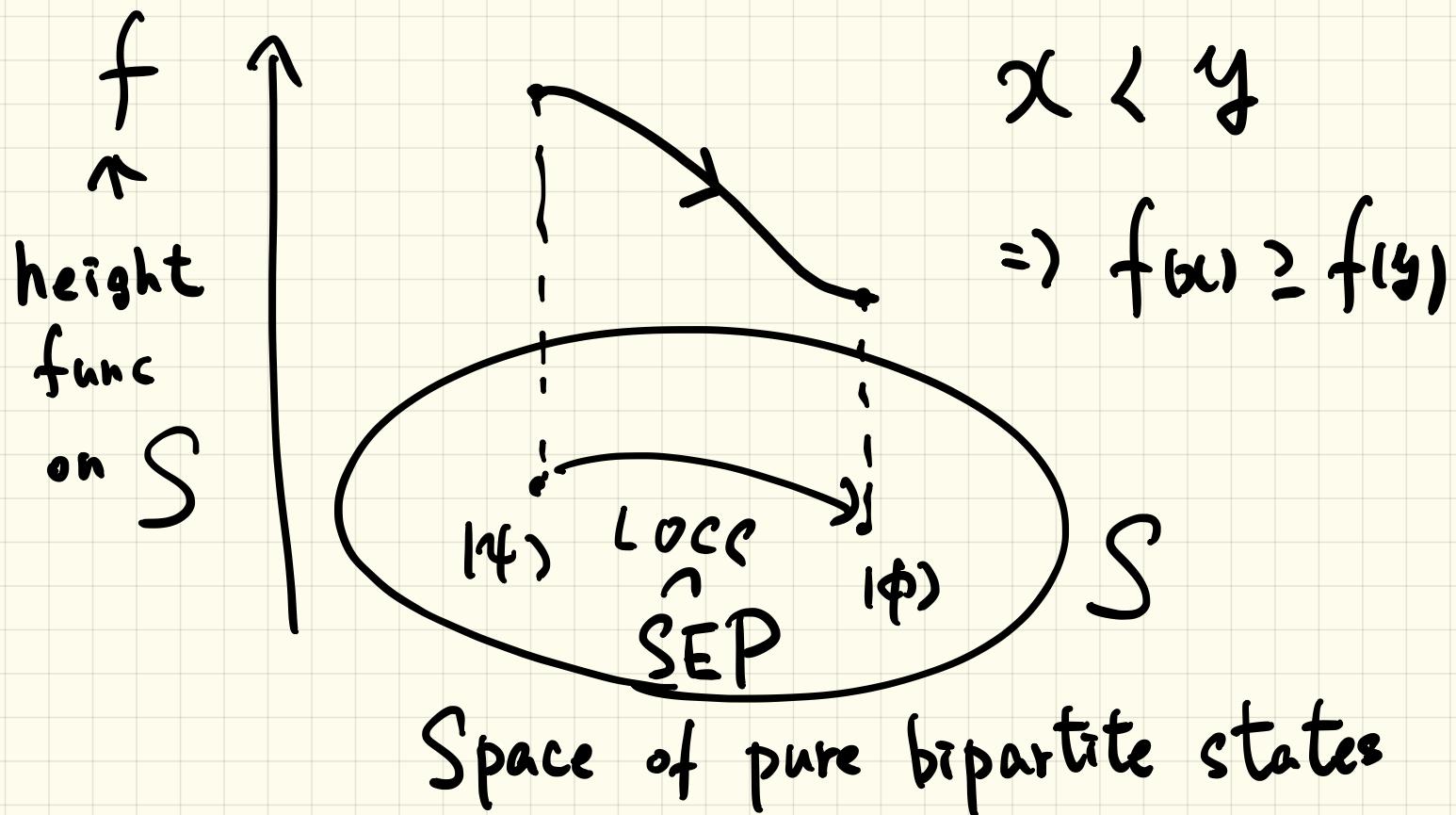
$|\psi \rangle$ is transformed to $|\phi \rangle$

under SEP iff $\lambda^+ < \lambda^\phi$

$|\psi \rangle \xrightarrow{\text{SEP}} |\phi \rangle$ iff $\lambda^+ < \lambda^\phi$

-

• How to quantify the order?



Schur-concave function

$$x \prec y \Rightarrow f(x) \geq f(y)$$

NB. $f(x) \geq f(y)$ does not necessarily mean $x \prec y$

but. it means $y \succ x$

Quantum information measures (entanglement)

A1. Monotonicity

$$|\psi\rangle \xrightarrow{\text{Locc}} |\phi\rangle \Rightarrow f(\rho_\psi) \geq f(\rho_\phi)$$

A2. Vanishing for sep. states

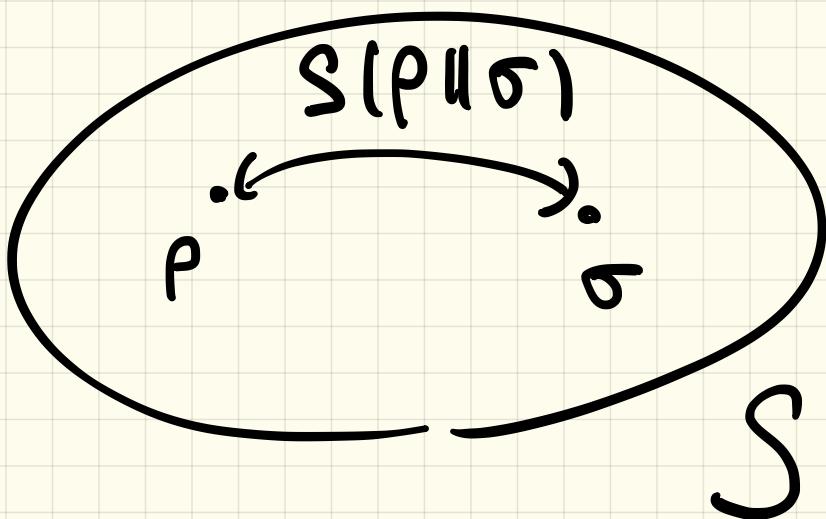
$$|\psi\rangle : \text{separable.} \Leftrightarrow f(\rho_\psi) = 0$$

A3. Convexity

$$f\left(\sum_i p_i \rho_i\right) \leq \sum_i p_i f(\rho_i)$$

Relative entropy

\approx "distance" bet quantum states



$$S(\rho \parallel \sigma) \quad \text{TP: } \text{Tr}(\rho) = 1 \Rightarrow \text{Tr}(h(\rho)) = 1$$

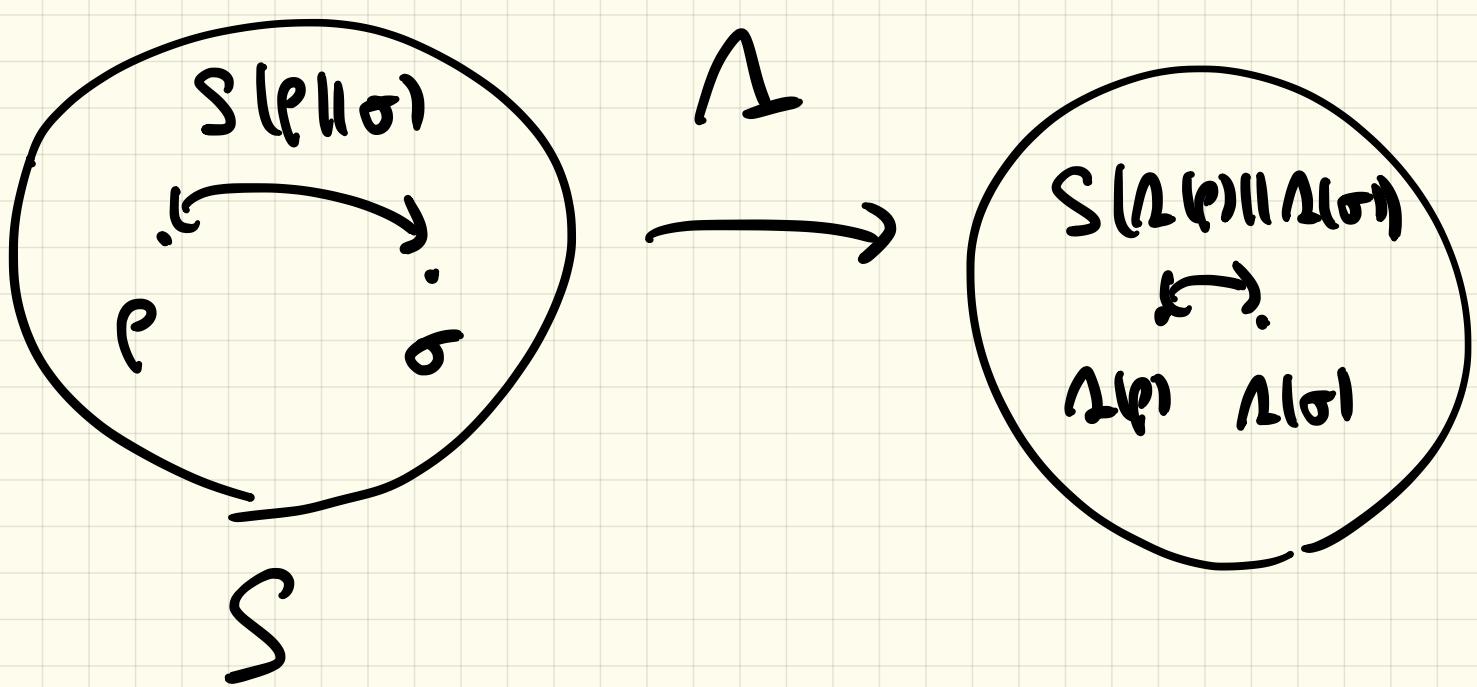
$$= \text{Tr}[\rho (\log \rho - \log \sigma)]$$

• Monotonicity Λ : SEP

$$S(\Lambda(\rho) \parallel \Lambda(\sigma))$$

$$\leq S(\rho \parallel \sigma)$$

C CPTP
 \uparrow
 completely
 positive
 trace preserving



- Positivity

$$S(\rho \parallel \sigma) \geq 0$$

- $S(\rho \parallel \sigma) = 0$

$$\Leftrightarrow \rho = \sigma$$

Example of CPTP

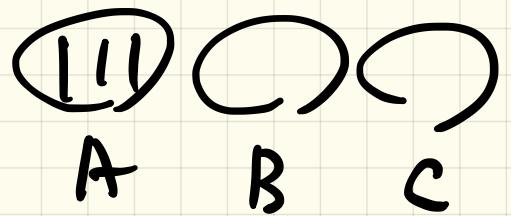
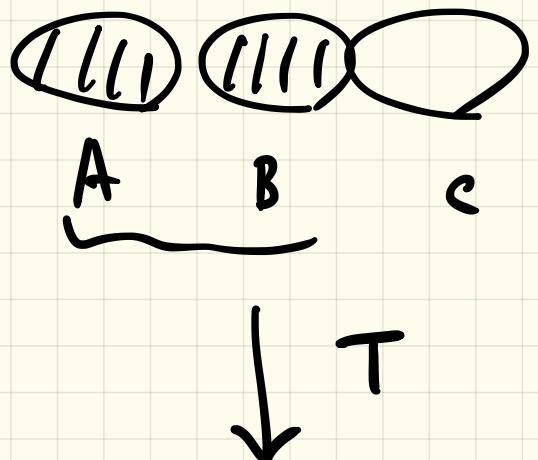
- Partial trace

suppose $\rho_{A \cup B}$

$$T = \mathbb{1}_A \otimes \text{Tr}_B$$

$$T : \rho_{A \cup B} \rightarrow \rho_A$$

$$\Rightarrow S(\rho_A \parallel \sigma_A) \leq S(\rho_{A \cup B} \parallel \sigma_{A \cup B})$$



Entanglement measures

- Entanglement entropy

$$\rho_A = \text{Tr}_{\bar{A}}(\rho)$$

$$S_A(\rho) \equiv -\text{Tr}_A[\rho_A \log \rho_A]$$

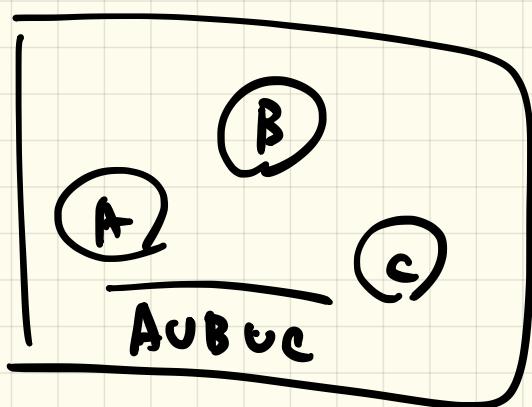
$$= \log d_A - S(\underbrace{\rho_A \parallel \mathbb{I}_A / d_A}_{\uparrow})$$

maximally
entangled
state

- From the monotonicity of rel. ent. (MONO)

Strong subadditivity

$$\begin{aligned} S_{A \cup C} + S_{B \cup C} & \stackrel{\text{SSA}}{\geq} S_A + S_B \end{aligned}$$



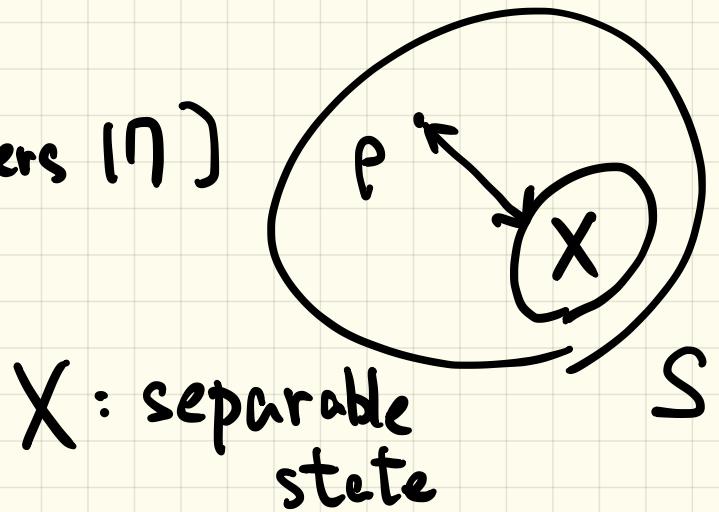
$\boxed{SSA \Leftrightarrow MONO}$

E.g. 2) Relative ent. of entanglement

$$E_R^X(\rho) \equiv \inf_{\sigma \in X} S(\rho \parallel \sigma)$$

In QFT.

[Holland - Saunders 17]



X: separable state

Q I T

separable states

CPTP

entangled states

entanglement
measures

Resource thy

free states

free operations

resource

resource measure

Ref)

- M. Nielsen

"An introduction to majorization
and its application to quantum mechanics"

- Plenio - Virmani quant-ph/0504163
- Chitambar - Gour 1806.06107

2. Quantum entanglement in QFT

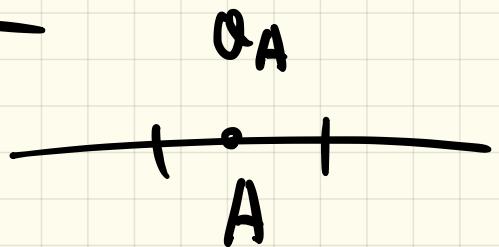
- Algebraic formulation

- Path-integral formulation

useful for proving inequalities
practical computation

Algebraic formulation

$$\mathcal{H} = \mathcal{H}_A \otimes \mathcal{H}_{\bar{A}}$$



• Reduced density matrix

$$\rho_A = \text{Tr}_{\bar{A}}[\rho]$$

$$\mathcal{O}_A = \mathcal{O}_A \otimes \mathbb{1}_{\bar{A}}$$

$$\langle \mathcal{O}_A \rangle = \text{Tr} [\rho (\mathcal{O}_A \otimes \mathbb{1}_{\bar{A}})]$$

$$= \text{Tr}_A [\rho_A \mathcal{O}_A] \quad \text{for any } \mathcal{O}_A$$

(Ω_A) invariant under transformation

$$\Omega_A \rightarrow U \Omega_A U^\dagger \quad U: \mathcal{H}_A \rightarrow \mathcal{H}_A$$

$$P_A \rightarrow U P_A U^\dagger$$

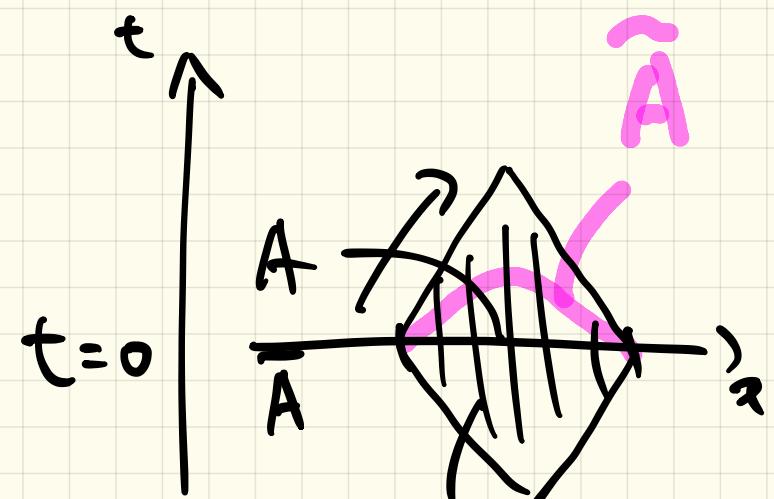
$\tilde{\Sigma}_A(P)$ is also inv.

U = time evolution
on A

If $D[\hat{A}] = D[A]$

then $\exists U$ s.t.

$$P_{\hat{A}} = U P_A U^\dagger$$



$D[A]$
domain of
dependence

AQFT

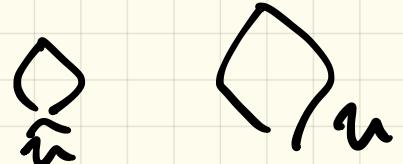
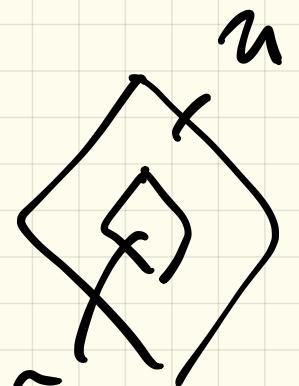
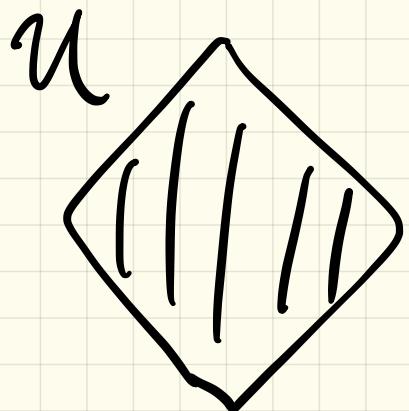
- Operator algebra \mathcal{A}_U associated with a spacetime region U

Postulates

$$1. \hat{U} \subset U$$

$$\Rightarrow \mathcal{A}_{\hat{U}} \subset \mathcal{A}_U$$

\rightsquigarrow Partial order



2. If u, \tilde{u} are spacelike

$$[\alpha_u, \alpha_{\tilde{u}}] = 0$$



3. Covariance [Haag's textbook]

Classification of α

- Type I (QMI)
- Type II
- Type III (QFT)
 α_{III_1}

Relative entropy in QFT

For a spacetime region

$$\mathcal{N} = D[A]$$

the rel. ent. between $| \psi \rangle, |\phi \rangle \in \mathcal{H}$

$$S_A(\rho_\psi \parallel \rho_\phi) = -\langle \psi | \log \Delta_{\psi|\phi} | \psi \rangle$$

$$\Delta_{\psi|\phi} = \rho_{\phi,A} \otimes \rho_{\psi,\bar{A}}^{-1} \quad [\text{Araki 76}]$$

(relative modular operator)

$$\rho_\psi = |\psi\rangle\langle\psi|$$

$$\rho_\phi = |\phi\rangle\langle\phi|$$

$$S_A(\rho_{\gamma} \parallel \rho_{\phi}) = S(\rho_{\gamma,A} \parallel \rho_{\phi,A})$$

Proof

$$\begin{aligned}
 S_A(\rho_{\gamma} \parallel \rho_{\phi}) &= -\text{tr}[\rho_{\gamma} \log \Delta_{\gamma(\phi)}] \\
 &= -\text{tr}[\rho_{\gamma} (\log \rho_{\phi,A} \otimes \mathbb{1}_{\bar{A}} \\
 &\quad - \mathbb{1}_A \otimes \log \rho_{\gamma,\bar{A}})] \\
 &= -\text{tr}_A[\rho_{\gamma,A} \log \rho_{\phi,A}] \quad \text{Schmidt} \\
 &\quad + \text{tr}_{\bar{A}}[\rho_{\gamma,\bar{A}} \log \rho_{\gamma,\bar{A}}] \quad \text{decomp.} \xrightarrow{(\bar{A} \rightarrow A)}
 \end{aligned}$$

$$\begin{aligned}
 &= -\text{tr}_A [P_{\Phi,A} \log P_{\Phi,A}] \\
 &\quad + \text{tr}_A [P_{\Psi,A} \log P_{\Psi,A}] \\
 &= \text{tr}_A [P_{\Psi,A} (\log P_{\Psi,A} - \log P_{\Phi,A})] \\
 &= S(P_{\Psi,A} \parallel P_{\Phi,A})
 \end{aligned}$$

//
This is the case for Type I

The new definition of rel. ent
is valid for Type III (w/ Tomita-Takesaki)
theory

Ref)

• Witten 1803.04993